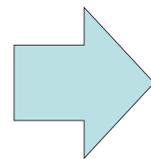


Static Analysis Combining Logical and Algebraic Abstractions

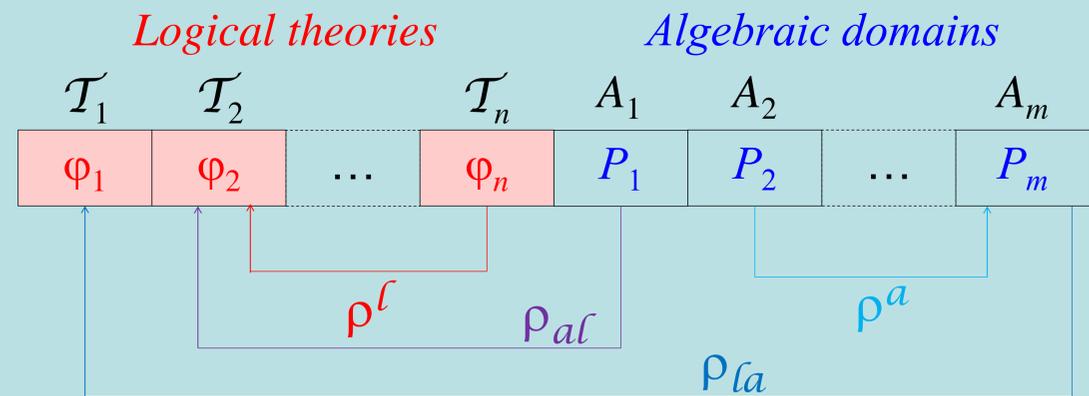
Proof theoretic/logical abstractions are a particular case of algebraic abstractions.

Nelson-Oppen procedure is a particular case of reduced product.

A new combination of logical and algebraic abstractions can be proposed by this unifying point of view.



Reduced Product of Logical and Algebraic Domains



Easy introducing new abstractions on either side
 \Rightarrow Extensible expressive static analyzers / verifiers.

SMT Solver Based Logical Abstract Domains

Abstract Properties

Universal representation of abstract properties by a set \mathcal{A} of quantifier-free first-order logical formulae.
 Each domain contains only one theory \mathcal{T} which may be decidable, deductive and complete on \mathcal{A} .

Abstract Operators

- ❖ Top: $\top \Rightarrow \text{tt}$
- ❖ Bottom: $\perp \Rightarrow \text{ff}$
- ❖ Abstract Intersection: $\prod \Rightarrow \bigwedge$
- ❖ Abstract Union: $\sqcup \Rightarrow \bigvee$

Abstract Implication

SMT solver is used for the abstract orders \sqsubseteq . We ask SMT solver whether $(\exists x_{\psi_1} \cup x_{\psi_2} :) \psi_1 \wedge \neg \psi_2$ is satisfied.
 ➤ Yes $\Rightarrow \psi_1 \not\sqsubseteq \psi_2$
 ➤ No $\Rightarrow \psi_1 \sqsubseteq \psi_2$

Abstractions

For a given logical abstract domain, we need to design an abstraction algorithm $\alpha \in F(x, f, p) \rightarrow \mathcal{A}$ to abstract properties in \mathcal{A} .

- Quantifier elimination
- Literal elimination

Abstract Transformers

- ✓ Invertible assignment:
 $f_a[[x:=t]]\psi \triangleq \psi[x \leftarrow t_x^{-1}]$
- ✓ Non-invertible assignment:
 $f_a[[x:=t]]\psi \triangleq \exists x' : \psi[x/x'] \wedge x = t$
- ✓ Test: $p[[\varphi]]\psi \triangleq \psi \wedge \varphi$

Widening

- ❑ Designing a universal widening for logical abstract domain is difficult.
- ❑ We can always ask the end-user.
- ❑ Alternative method: widening in polyhedral abstract domain.